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2. Because same layout between processes can protect the unity of the shared memory. Same structure assures that each process will visit all the elements of Shm in a same order.
3. memcpy() decreases the number of visiting memory, because memcpy() copy data by page size, but element-wise assignments copy data by bytes. Thus, memcpy() is more efficient.

[leiquanpan@shell studio13]$ ./leader

data list is of 0 element is: 19 for local, 19 for memory region

data list is of 1 element is: 92 for local, 92 for memory region

data list is of 2 element is: 81 for local, 81 for memory region

data list is of 3 element is: 8 for local, 8 for memory region

data list is of 4 element is: 65 for local, 65 for memory region

data list is of 5 element is: 63 for local, 63 for memory region

data list is of 6 element is: 1 for local, 1 for memory region

data list is of 7 element is: 76 for local, 76 for memory region

data list is of 8 element is: 10 for local, 10 for memory region

data list is of 9 element is: 21 for local, 21 for memory region

[leiquanpan@shell studio13]$ ./follower

data list is of 0, 19 for memory region

data list is of 1, 92 for memory region

data list is of 2, 81 for memory region

data list is of 3, 8 for memory region

data list is of 4, 65 for memory region

data list is of 5, 63 for memory region

data list is of 6, 1 for memory region

data list is of 7, 76 for memory region

data list is of 8, 10 for memory region

data list is of 9, 21 for memory region

1. Because shared memory has already been cerated by leader. ftruncate() may change the size of the shared memory, which may cause big problem to other process who shared this memory.
2. Because using three seperate guards assure that the leader can receive the right singal from the follower, so it won’t be a deadlock hazard. Since the guard’s usage is like a barrier to control leader write and follower read at a right order. Therefore, I don’t think there will be any data races.
3. Therefore, its bandwith is 3.17 GB/S

size:1000000

[leiquanpan@shell studio13]$ time ./leader

real 0m0.295s

user 0m0.284s

sys 0m0.010s